I will talk about low linear orderings with computable presentation. An $X$-computable linear ordering is called low, if $X' \leq_T \emptyset'$. C.G. Jockusch and R.I. Soare [5] proved that any noncomputable c.e. degree contains linear ordering with no computable presentation. Therefore, there exists a low linear ordering with no computable copy.

R.G. Downey, M.F. Moses [2] proved that any low discrete linear ordering has a computable copy (a linear ordering is called discrete, if any element has both a successor and a predecessor). It is a natural to ask (R.G. Downey, [1]) — is there a property $P$ of order types which guarantees that if $L$ is low and $P(L)$ then $L$ has a computable presentation?

The author [4] proved that any low strongly $\eta$-like linear ordering is isomorphic to a computable one (a linear ordering $L$ is called strongly $\eta$-like, if $L \cong \sum_{q \in Q} f(q)$, where $|\text{rang}(f)| < +\infty$). Also the author showed that any low 1-quasidiscrete has a computable copy.

**Definition** A linear ordering is called $k$-quasidiscrete, if any equivalence class either is infinite or contains at most $k$ elements, where $x \sim y$ iff there are only finite set of $z$ such that $x \leq_L z \leq_L y$ or $y \leq_L z \leq_L x$.

**Theorem** Any low $k$-quasidiscrete linear ordering is a computable presentable ordering.

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