

Decidable Boolean algebras of elementary characteristics (1,0,1)

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Definition

A computable Boolean algebra is said to be n -constructive if there exists an algorithm determining whether the given tuple satisfies given finite Σ_n -formula.

Definition

A strongly constructive Boolean algebra is one for which such an algorithm exists for all formulas of the predicate calculus.

Definition

Boolean algebra is decidable if there exists its strongly constructive isomorphic copy.

a — **atom**, if $a \neq 0$ and $\forall b(b < a \rightarrow b = 0)$.

$\mathbf{At}(\mathfrak{A})$ — set of atoms of \mathfrak{A} .

a — **atomic element**, if $\forall x \leq a (x \neq 0 \rightarrow (\exists y \leq x (y \in \mathbf{At}(\mathfrak{A}))))$.

$\mathbf{Atm}(\mathfrak{A})$ — ideal of atomic elements of \mathfrak{A} .

a — **atomless element**, if $\forall x \leq a (x \notin \mathbf{At}(\mathfrak{A}))$.

$\mathbf{Als}(\mathfrak{A})$ — ideal of atomless elements of \mathfrak{A} .

$\mathbf{E}(\mathfrak{A}) = \mathbf{Atm}(\mathfrak{A}) + \mathbf{Als}(\mathfrak{A})$ — Ershov-Tarski ideal.

$At(\mathfrak{A}), Als(\mathfrak{A}), Atm(\mathfrak{A}), E(\mathfrak{A})$

$ch(\mathfrak{A}) = (1, 0, 1) \Leftrightarrow \mathfrak{A}/E$ — nontrivial atomless Boolean algebra.

$$S \subseteq \{At(\mathfrak{A}), Als(\mathfrak{A}), Atm(\mathfrak{A}), E(\mathfrak{A})\}$$

Is \mathfrak{A} decidable if predicates from S are computable?

$$S = \{At(\mathfrak{A}), Als(\mathfrak{A})\} \Rightarrow \text{yes!}$$

$$S = \{At(\mathfrak{A})\} \Rightarrow \text{no}$$

$$S = \{Als(\mathfrak{A}), Atm(\mathfrak{A}), E(\mathfrak{A})\} \Rightarrow \text{no}$$

Unconsidered cases:

$$S = \{At(\mathfrak{A}), Atm(\mathfrak{A})\} \Rightarrow ?$$

$$S = \{At(\mathfrak{A}), E(\mathfrak{A})\} \Rightarrow ?$$

Theorem

Let \mathfrak{A} be a computable Boolean algebra of elementary characteristics $(1, 0, 1)$. If $At(\mathfrak{A})$ and $Atm(\mathfrak{A})$ are computable then \mathfrak{A} is decidable.

Theorem

There exists a computable Boolean algebra \mathfrak{A} of elementary characteristics $(1, 0, 1)$ with computable $At(\mathfrak{A})$ and $E(\mathfrak{A})$ which is not decidable.

Description of Δ_6^0 -computable Boolean algebras

Let $T = (\text{Atm} \rightarrow F) + \text{Atm}$,

where $\text{Atm} \rightarrow F = \{x \mid \forall z \leq x (z \in F \vee z \notin \text{Atm})\}$

F — Frechet ideal.

Theorem

\mathfrak{A} is Δ_6^0 -computable Boolean algebra if and only if there exists a computable Boolean algebra \mathfrak{C} such that $\mathfrak{C}/T \cong \mathfrak{A}$.